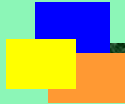


國立清華大學 電機工程學系
EE2410 Data Structure

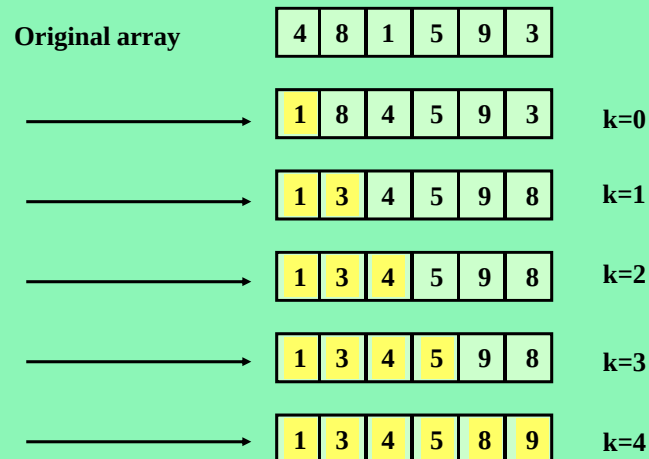


Chapter 3 Stacks and Queues

Outline

- ➡ • **Templates in C++**
 - **The Stack Abstract Data Type**
 - **The Queue Abstract Data Type**
 - **SubTyping and Inheritance in C++**
 - **A Mazing Problem**
 - **Evaluation of Expressions**

Example of Selection Sort



ch3-3

Selection Sort Using Template

```

1.  template <class KeyType>
2.  void sort (KeyType *a, int n)
3.  // sort the n KeyType a[0] to a[n-1] into nondecreasing order
4.  {
5.      for(int k=0; k < n; k++){
6.          int smallest = k;
7.          // find the smallest KeyType in a[k] to a[n-1]
8.          for (int j = k+1; j < n; j++){
9.              if(a[j] < a[smallest]) smallest = j;
10.         }
11.         KeyType temp = a[k]; a[k] = a[smallest]; a[smallest] = temp;
12.     }
13. }
14. main(){
15.     float real_array[100];
16.     int int_array[250];
17.     ... // assume that the arrays have been initialized
18.     sort(real_array, 100);
19.     sort(int_array, 250);
20.     ...
21. }
```

ch3-4

Concept of Template

- **A Template**
 - also called **parameterized types**
 - It may be viewed as a **variable** which can be **instantiated** to any data type when a function or class is used
- **Copy Constructor**
 - is a special constructor which is invoked when **an object is initialized with another object**
 - e.g., **KeyType temp=a[k];**
 - might be overloaded if the parameterized type is not a pre-defined data type in C++

ch3-5

Container Class

```
1. class Bag
2. {
3. public:
4.     Bag (int MaxSize = DefaultSize); // constructor
5.     ~Bag(); // destructor
6.     void Add(int); // insert an integer into Bag
7.     void Delete(int &); // delete an integer from Bag
8.     Boolean IsFull(); // return TRUE if the bag is full; FALSE otherwise
9.     Boolean IsEmpty(); // return TRUE if the Bag is empty; FALSE otherwise
10. private:
11.     void Full(); // action when bag is full
12.     void Empty(); // action when the bag is empty
13.
14.     int *array;
15.     int MaxSize; // size of array
16.     int top; // highest position in array that contains an element
17. };
```

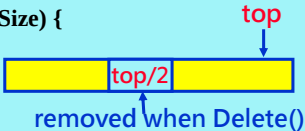
ch3-6

Implementation of Bag

```

1.  Bag::Bag(int MaxBagSize): MaxSize (MaxBagSize) {
2.      array = new int[MaxSize]; top = -1;
3.  }
4.  Bag::~~Bag() { delete [] array; }
5.  inline Boolean Bag::IsFull(){
6.      if(top == MaxSize - 1) return TRUE; else return FALSE;
7.  }
8.  inline Boolean Bag::IsEmpty(){
9.      if(top == -1) return TRUE; else return FALSE;
10. }
11. inline void Bag::Full(){ cout << "Bag is full" << endl; }
12. inline void Bag::Empty(){ cout << "Bag is empty" << endl; }
13. void Bag::Add(int x) { if(IsFull()) Full; else array[++top] = x; }
14. int *Bag::Delete(int &x){
15.     if( IsEmpty()) { Empty(); return(0); }
16.     int deletePos = top/2; x = array[deletePos];
17.     for(int k=deletePos; k<top; k++){
18.         array[k] = array[k+1]; // compact upper half of array
19.     }
20.     top--; return(&x);

```



ch3-7

Parameterized Container Class

```

1.  template <class Type>
2.  class Bag
3.  {
4.  public:
5.      Bag (int MaxSize = DefaultSize); // constructor
6.      ~Bag(); // destructor
7.      void Add(const Type&); // insert an element into Bag
8.      Type *Delete(Type&); // delete middle element from Bag
9.      Boolean IsFull(); // return TRUE if the bag is full; FALSE otherwise
10.     Boolean IsEmpty(); // return TRUE if the Bag is empty; FALSE otherwise
11. private:
12.     void Full(); // action when bag is full
13.     void Empty(); // action when the bag is empty
14.
15.     Type *array;
16.     int MaxSize; // size of array
17.     int top; // highest position in array that contains an element
18. };
19. main(){
20.     Bag<int> a; Bag<Rectangle> r; ...
21. }

```

end-8

Outline

- Templates in C++
- ➡ • **The Stack Abstract Data Type**
- The Queue Abstract Data Type
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- A Mazing Problem
- Evaluation of Expressions

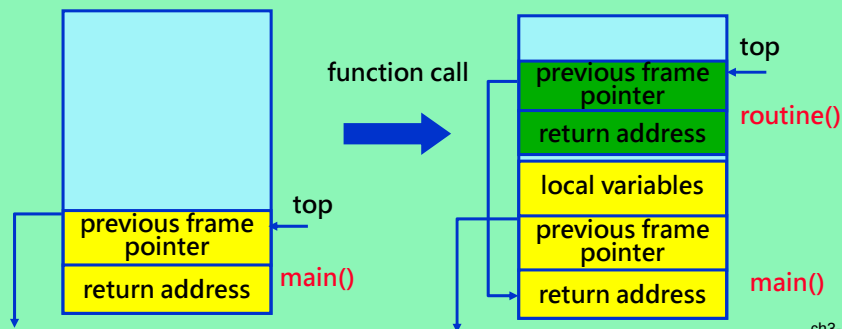
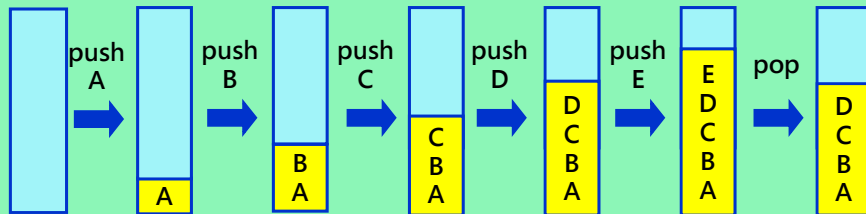
ch3-9

Stack

- **Ordered List**
 - Suppose $A = a_0, a_1, \dots, a_{n-1}$, where $n \geq 0$
 - a_i is called an **atom**, or an **element**
- **Stack: Last-In First-Out (LIFO)**
 - is a special case of ordered list
 - the **additions** and **deletions** are made at one end called the **top**
 - E.g., Given a stack $S = (a_0, a_1, \dots, a_{n-1})$
 - a_0 is the bottom element
 - a_{n-1} is the top element

ch3-10

Example: Stack Operation



ch3-11

ADT of Stack

```

1. template <class KeyType>
2. class Stack
3. {
4. // objects: A finite ordered list of zero or more elements
5. public:
6.     Stack (int MaxStackSize = DefaultSize);
7.     Boolean IsFull(); // return TRUE if Stack is full; FALSE otherwise
8.     Boolean IsEmpty(); // return TRUE if Stack is empty; FALSE otherwise
9.     void Add(const KeyType&);
10.        // if IsFull(), return 0;
11.        // else insert an element to the top of the Stack
12.     KeyType *Delete(KeyType&);
13.        // if IsEmpty(), then return 0;
14.        // else remove and return a pointer to the top element
15. private:
16.     int top;
17.     KeyType *stack;
18.     int MaxSize;
19. };
    
```

ch3-12

Member Functions of Stack

```
1. template <class KeyType>
2. Stack<KeyType>::Stack (int MaxStackSize) : MaxSize (MaxStackSize)
3. {
4.     stack = new KeyType[MaxSize];
5.     top = -1;
6. }
7.
8. template <class KeyType>
9. inline Boolean Stack<KeyType>::IsFull()
10. { if(top == MaxSize - 1) return TRUE; else return FALSE; }
11.
12. template <class KeyType>
13. inline Boolean Stack<KeyType>::IsEmpty()
14. { if (top == -1) return TRUE; else return FALSE; }
15.
```

ch3-13

Push & Pop of Stack

```
1. template <class KeyType>
2. void Stack<KeyType>::Add(const KeyType& x) // the push operation
3. {
4.     if( IsFull() ) StackFull();
5.     else {
6.         stack[++top] = x;
7.     }
8. }
9.
10. template <class KeyType>
11. KeyType* Stack<KeyType>::Delete(KeyType& x) // the pop operation
12. {
13.     if (IsEmpty() ) StackEmpty();
14.     else {
15.         x = stack[top--]; // assigned the item being deleted to x
16.         return( &x );
17.     }
18. }
```

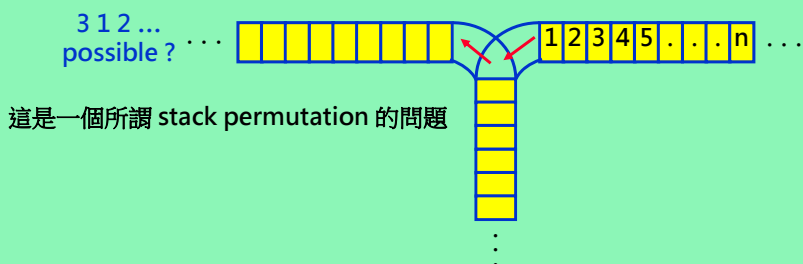


The diagram shows two states of a stack. In the first state (push), the stack contains elements A, B, C, D, E, and a new element 'x' is being added at the top. The 'top' pointer points to 'x'. In the second state (pop), the element 'x' has been removed, and the 'top' pointer now points to 'E'.

ch3-14

Railroad Switching System

- **Switching Rule**
 - **Initial:** train 1, 2, ..., n in the top right track segment
 - **Movement:**
 - (1) from **top-right** to the **vertical segment** one at a time
 - (2) from the **vertical** to the **top-left** segment one at a time
 - (3) The vertical segment operates like a **stack**
- **Question: What output permutations are not possible?**



ch3-15

Outline

- Templates in C++
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- ➡ • **The Queue Abstract Data Type**
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ch3-16

Queue

- **A Queue**

- is an ordered list in which all **insertions** take place at one end and all **deletions** take place at the opposite end
- is also called **First-In First-Out (FIFO)**

- **Example**

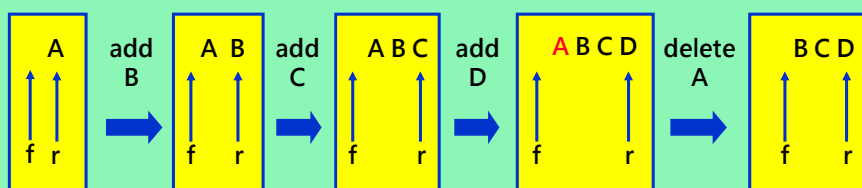
- Given a queue $Q = (a_0, a_1, \dots, a_{n-1})$
- a_0 is the **front** element, a_{n-1} is the **rear** element
- a_i is **behind** a_{i-1} for $1 \leq i \leq n$

ch3-17

Data and Operations of Queue

```
1. template <class KeyType>
2. class Queue {
3. public: ...
4. private:
5.     int front;           // front: index of the first element to be retrieved
6.     int rear;            // rear: index of the last element
7.     KeyType *queue;      // KeyType array
8.     int MaxSize;
9. };
```

f = r = -1 initially

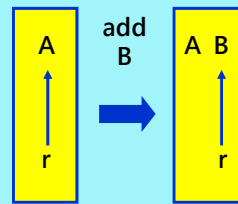


ch3-18

Member Functions of Queue

```

1.  template <class KeyType>
2.  Queue<KeyType>::Queue (int MaxQueueSize) : MaxSize (MaxQueueSize)
3.  {
4.      queue = new KeyType[MaxSize];
5.      front = rear = -1;
6.  }
7.
8.  template <class KeyType>
9.  inline Boolean Queue<KeyType>::IsFull()
10. {
11.     if( rear == MaxSize -1 ) return TRUE;
12.     else return FALSE;
13. }
14.
15. template <class KeyType>
16. inline Boolean Queue<KeyType>::IsEmpty()
17. {
18.     if( front == rear ) return TRUE;
19.     else return FALSE;
20. }
    
```

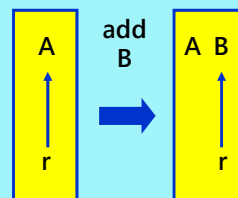


ch3-19

Add and Delete for Queue

```

1.  template <class KeyType>
2.  void Queue<KeyType>::Add (const KeyType& x)
3.  // add x to the queue
4.  {
5.      if( IsFull() ) QueueFull();
6.      else queue[++rear] = x;
7.  }
8.
9.  template <class KeyType>
10. KeyType *Queue<KeyType>::Delete(KeyType& x)
11. // remove front element from the queue
12. {
13.     if( IsEmpty() ) { QueueEmpty(); return(); }
14.     x = queue[++front];
15.     return x;
16. }
    
```



ch3-20

Example: Job Scheduling

- **In Operating System**

- jobs are processed in the order they enter the system
if no priority is set on jobs

front	rear	Q[0]	Q[1]	Q[2]	Q[3]	Q[4]	Q[5]	comments
-1	-1							queue is empty
-1	0	J1						Job 1 joins Q
-1	1	J1	J2					Job 2 joins Q
-1	2	J1	J2	J3				Job 3 joins Q
0	2		J2	J3				Job 1 leaves Q
0	3		J2	J3	J4			Job 4 joins Q
1	3			J3	J4			Job 2 leaves Q

ch3-21

Worst-Case Scenario

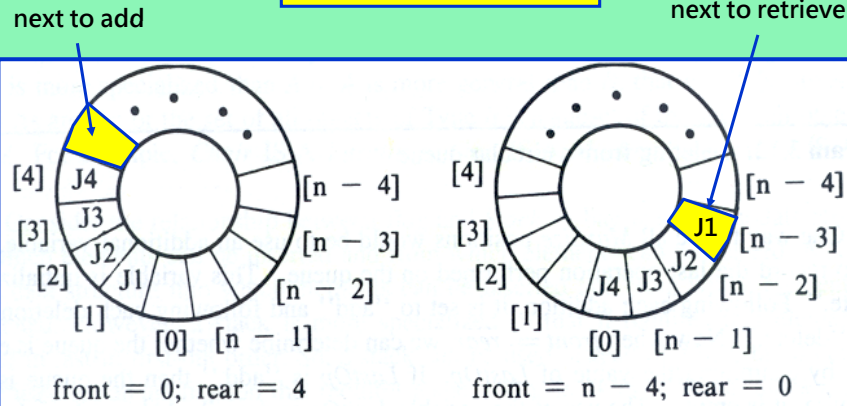
front	rear	Q[0]	Q[1]	Q[2]	...	Q[n-1]	Next Operation
-1	n-1	J1	J2	J3	...	J _n	initial state
0	n-1		J2	J3	...	J _n	delete J1
-1	n-1	J2	J3	J4	...	J _{n+1}	add J _{n+1} (J2 to J _n are moved)
0	n-1		J3	J4	...	J _{n+1}	delete J2
-1	n-1	J3	J4	J5	...	J _{n+2}	add J _{n+2}

In the above job scheduling,
it takes **n-1 steps** to add a new job

ch3-22

Circular Queue

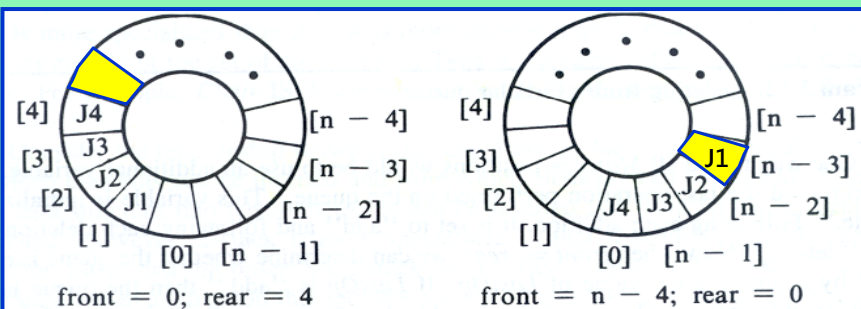
Queue size: n
Jobs: J1, J2, J3, J4



ch3-23

Add An Element to Circular Q

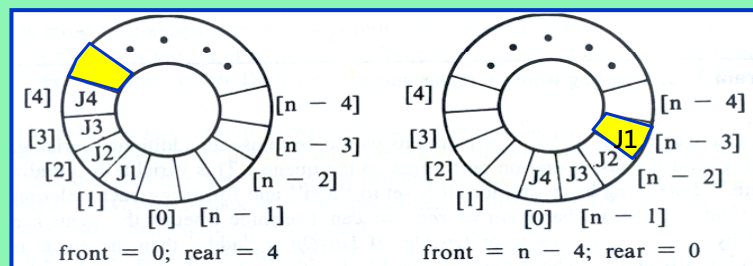
```
template <class KeyType>
void Queue<KeyType>::Add (const KeyType& x)
{
    int new_rear=(rear+1)% MaxSize;
    if( front == new_rear) QueueFull();
    else queue[rear == new_rear] = x;
}
```



ch3-24

Delete An Element From Circular Q

```
template <class KeyType>
KeyType *Queue<KeyType>::Delete (KeyType& x)
// remove front element from queue
{
    if ( front == rear ) { QueueEmpty(); return(0); }
    front = (front + 1) % MaxSize;
    x = queue[front]; return (&x);
}
```



ch3-25

DeQue

- **Definition**

- A double-ended queue (Deque) is a linear list in which additions and deletions may be made at either end

- **Exercises**

- Design a data representation that maps a deque into a one-dimensional array
- Write algorithms to add and delete elements from either end of the queue

ch3-26

Outline

- Templates in C++
- The Stack Abstract Data Type
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- ➡ • **SubTyping and Inheritance in C++**
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ch3-27

SubType and Inheritance

- **IS-A Relation**
 - Chair IS-A furniture, Lion IS-A Mammal
 - Rectangle IS-A Polygon
 - Stack IS-A Bag
- **Inheritance**
 - Is used to express IS-A relationship between ADTs.
 - **derived class** IS-A **base class**
 - C++ provides a mechanism called **public inheritance**
 - Ex: **class Stack: public Bag**
 - The **inherited members** have the same level of **access** in the derived class as in the base class

ch3-28

Bag and Stack

```
class Bag
{
public:
    Bag (int MaxSize = DefaultSize); // constructor
    ~Bag(); // destructor
    virtual void Add(int);
    virtual int *Delete(int&);
    virtual Boolean IsFull();
    virtual Boolean IsEmpty();
protected:
    virtual void Full();
    virtual void Empty();
    int *array; int MaxSize; int top; // data members
};
```

- (1) Interface of Bag will be reused in Stack
- (2) The implementation of virtual functions may be re-defined in the derived class

```
class Stack : public Bag
{
public:
    Stack (int MaxSize = DefaultSize);
    ~Stack();
    int *Delete(int&); // delete the element from stack
};
```

constructor, destructor cannot be reused

ch3-29

Example of Derived Class

```
Stack::Stack(int MaxStackSize) : Bag(MaxStackSize) { }
// Constructor for Stack calls constructor for Bag

Stack::~~Stack(): ~Bag() { }
// Destructor for Bag is automatically called when Stack is destroyed.
// This ensures that array is deleted

int *Stack::Delete(int& x)
{
    if ( IsEmpty() ) { Empty(); return(0); }
    x = array[ top-- ];
    return(&x);
}
```

Example: Stack s(3), int x
 s.Add(1); s.Add(2); s.Add(3);
 // Stack::Add is not defined, so use Bag::Add instead
 s.Delete(x)
 // uses Stack::Delete, which calls Bag::IsEmpty and Bag::Empty
 because these have not been redefined in Stack

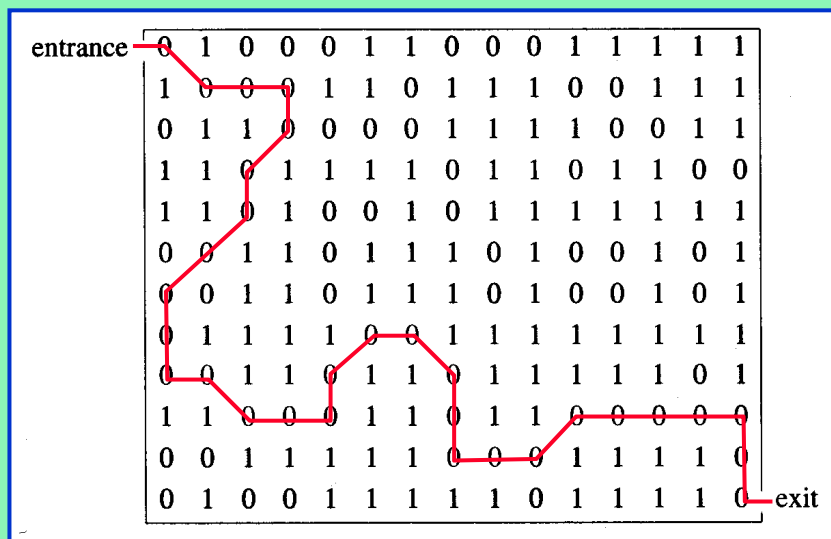
ch3-30

Outline

- Templates in C++
- The Stack Abstract Data Type
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ch3-31

An Example Maze

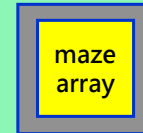


ch3-32

Data Structure

- **Maze**

- Is represented as a two-dimensional array `maze[i][j]`
- `maze[i][j]=0`: location that can be passed through
- `maze[i][j]=1`: blocked location
- **Entrance**: `maze[0][0]`
- **Exit**: `maze[m][p]`



To model the boundary

- **To model border condition**

- The array is declared as `maze[m+2][p+2]`
- I.e., the original maze array is surrounded by a border of ones

ch3-33

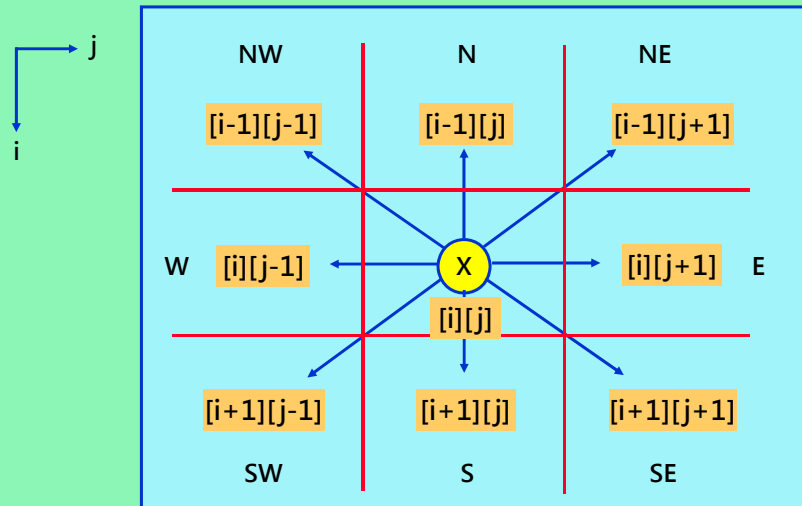
Strategy of Searching

- **As a rat walks through the maze**

- (1) He **picks a valid move** from the current position
 - E.g., starting from north and looking clockwise
- (2) Put the selected move into a **stack**
 - So that he can return from a **dead path**
- (3) He learns **not to make the same mistake twice**
 - Avoid getting into a cell visited before
 - A 2-dimensional array, `mark[m+2][p+2]` is used
 - The mark array records the cells visited before

ch3-34

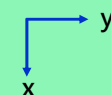
Allowable Moves



ch3-35

Coordinates of the Next Move

The coordinates of the next move is computed by the following data structure



```
struct offsets
{
    int x, y;
};
enum directions {
    N, NE, E, SE, S, SW, W, NW
};
offsets move[8];
```

→ The SW of (i, j) will be (g, h)
where
 $g = i + \text{move[SW].x};$
 $h = j + \text{move[SW].y}$

q	move[q].x	move[q].y
N	-1	0
NE	-1	1
E	0	1
SE	1	1
S	1	0
SW	1	-1
W	0	-1
NW	-1	-1

ch3-36

First Pass at Finding A Path Through a Maze

```

initialize stack to the maze entrance coordinates and direction east;
while (stack is not empty)
{
    (i,j,dir) = coordinates and direction deleted from top of stack ;
    while (there are more moves)
    {
        (g,h) = coordinates of next move ;
        if (( g == m ) && ( h == p )) success ;
        if ( (!maze [g ][h]) // legal move
            && (!mark [g ][h]) // haven't been here before
        {
            mark [g ][h] = 1 ;
            dir = next direction to try ;
            add (i,j,dir) to top of stack;
            i = g; j = h; dir = north;
        }
    }
}
cout << "no path found" << endl ;

```

An item of the stack:
 struct items {
 int *x, y, dir*;
 }

ch3-37

Example: A Mazing Problem

Row 0					
Row 1		0	0	0	
Row 2		0	1	1	
Row 3		1	0	0	
Row 4					
	C0	C1	C2	C3	C4

	stack
	top → (3, 2, E)
	(2, 1, SE)
	(1, 2, SW)
	(1, 1, E)

stack right before success

Current Position	Next Legal Move	Stack operation
(1, 1)	(1, 2, E)	Push (1, 1, E)
(1, 2)	(1, 3, E)	Push (1, 2, E)
(1, 3)	No legal move	Pop to backtrack
(1, 2)	(2, 1, SW)	Push (1, 2, SW)
(2, 1)	(3, 2, SE)	Push (2, 1, SE)
(3, 2)	(3, 3, E) success!	Pop out the entire stack

Complete path:

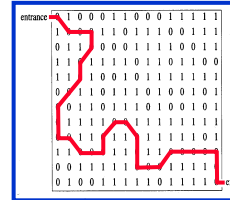
(3, 3) ← (3, 2, E) ← (2, 1, SE) ← (1, 2, SW) ← (1, 1, E)

ch3-38

```

void path(int m, int p)
// Output a path (if any) in the maze; maze[0][i] = maze [m + 1][i] =
// maze [j][0] = maze [j][p + 1] = 1, 0 ≤ i ≤ p + 1, 0 ≤ j ≤ m + 1.
{
    // start at (1,1)
    mark[1][1] = 1 ;
    Stack<items> stack(m*p) ;
    items temp ;
    temp.x = 1 ; temp.y = 1 ; temp.dir = E ;
    stack.Add(temp) ;
    while (!stack.IsEmpty()) // stack not empty
    {
        temp = *stack.Delete (temp) ; // unstack
        int i = temp.x ; int j = temp.y ; int d = temp.dir ;
        while (d < 8) // move forward
        {
            int g = i + move[d].a ; int h = j + move[d].b ;
            if ((g == m) && (h == p)) { // reached exit
                // output path
                cout << stack ;
                cout << i << " " << j << endl ; // last two squares on the path
                cout << m << " " << p << endl ;
                return;
            }
            if ((!maze[g][h]) && (!mark[g][h])) { // new position
                mark[g][h] = 1 ;
                temp.x = i ; temp.y = j ; temp.dir = d+1 ;
                stack.Add(temp) ; // stack it
                i = g ; j = h ; d = N ; // move to (g,h)
            }
            else d++ ; // try next direction
        }
    }
    cout << "no path in maze " << endl ;
}

```



9

Backward Search

Row 0					
Row 1		3	3	4	
Row 2		2	∞	∞	
Row 3		∞	1	0	
Row 4					
	C0	C1	C2	C3	C4

after step1

Row 0					
Row 1		3	3	4	
Row 2		2	∞	∞	
Row 3		∞	1	0	
Row 4					
	C0	C1	C2	C3	C4

after step2

Backward_search_algorithm

```

{
    Step 1:
        compute the distance to the destination
        of each node by wave-front propagation
    Step 2:
        find a shortest path from the source to
        the destination node by picking up the
        nodes with a shortest distance
}

```

ch3-40

Outline

- Templates in C++
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- ➡ • **Evaluation of Expressions**

ch3-41

Types of Expression

- **Arithmetic Expression**
 - For example: $X = A/B - (C + D * E - A * C)$
 - The evaluation of this expression is critical in enabling high-level programming
 - An expression consists of
 - (1) **Operands:** A, B, C, D, E
 - (2) **Operator:** plus, minus, times, and divide
 - (3) **Delimiter:** like parenthesis “(”, “)”
- **Boolean Expression**
 - The result in TRUE or FALSE
 - Use **relational** and **logical** operators
 - <, <=, ==, >, >=, &&, ||, !

ch3-42

Priority of Operator

- **Priority**

- The **order of operations** to be carried out in an expression
- Different order leads to **different results**

$$\begin{aligned} X &= A/B - C + D * E - A * C \\ &= ((4/2) - 2) + (3 * 3) - (4 * 2) \\ &= 0 + 9 - 8 \\ &= 1 \end{aligned}$$

$$\begin{aligned} X &= A / B - C + D * E - A * C \\ &= (4 / (2 - 2 + 3)) * (3 - 4) * 2 \\ &= (4/3) * (-1) * 2 \\ &= -2.6666 \end{aligned}$$

Times and divide
have higher priorities
→ Default

Plus and Minus
have higher priorities
→ Not default

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Priority of Operations in C++

Evaluation of **operators of the same priority** will proceed from **left to right**
E.g., $A/B * C \rightarrow (A/B) * C$

priority	operator
1	Unary minus, !
2	*, /, %
3	+, -
4	<, <=, >, >=
5	==, !=
6	&&
7	

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Postfix Notation

- **Compiler**
 - Translates an **expression** into a **sequence of machine codes**
 - It first re-writes the expression into a form called **postfix notation**
- **Infix notation**
 - The operators come **in-between** operands
- **Postfix notation**
 - The operators appear **after** its operands
- **Example**
 - Infix: $A / B - C + D * E - A * C$
 - Postfix: $A B / C - D E * + A C * -$

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Evaluation of Postfix Notation

- **Scanning the notation from left to right**
- **Store temporary result in T_i , $i \geq 1$**

Original expression: $AB/C-DE*+AC*-$	
Operation	Postfix
$T_1 = A/B$	$T_1 C - DE * + AC * -$
$T_2 = T_1 - C$	$T_2 DE * + AC * -$
$T_3 = D * E$	$T_2 T_3 + AC * -$
$T_4 = T_2 + T_3$	$T_4 AC * -$
$T_5 = A * C$	$T_4 T_5 -$
$T_6 = T_4 - T_5$	T_6

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The Virtues of Postfix Notation

- **Evaluation is easier on postfix notation**
 - The need for **parenthesis** is eliminated
 - The **priority** of operations is no longer relevant
 - Evaluation of each operator
 - (1) **Pop** correct number of **operands** from the stack
 - (2) Perform the **operation**
 - (3) **Push** the results onto the stack

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Evaluation Algorithm

```
void evaluation(expression e)
// Evaluate the postfix expression e. It is assumed that the last
// token ( a token is either an operator, operand, or '#' ) in e is '#'
// A function NextToken is used to get the next token from e.
{
    Stack<token> stack; // initialize the stack
    for( token x = NextToken(e); x != '#'; x = NextToken(e) ){
        if( x is an operand ) stack.Add(x) // add to stack
        else { // operator
            pop-up correct number of operands for operator x;
            perform the operation x and store the result onto
            the stack;
        }
    }
}
```

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Conceptual Infix to Postfix

- Algorithm**

- (1) Fully **parenthesize** the expression
- (2) **Move all operators** so that they replace their corresponding right parentheses
- (3) **Delete all parentheses**

Example: $A / B - C + D * E - A * C$

→ $((((A / B) - C) + (D * E)) - (A * C))$

→ $AB / C - DE * + AC * -$

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Ex1: From Infix to Postfix

- Translate $A+B*C$ to $ABC*+$**

Next token	Stack	Output
None	Empty	None
A	Empty	A
+	+	A
B	+	AB
*	+ *	AB
The operator * has a higher priority than +, so it is placed on top of +		
C	+ *	ABC

- (1) All operands go directly to the output
 (2) When the input tokens are exhausted
 → All **operators** in the stack will be taken off

ch3-50

Ex2: From Infix to Postfix

- Translate $A*(B+C)/D$ to $ABC+*D/$

Next token	Stack	Output
None	Empty	None
A	Empty	A
*	*	A
(	*(	A
B	*(	AB
+	*(+	AB
C	*(+	ABC
)	*	ABC+
/	/	ABC+*
D	/	ABC+*D
'#' (ending character)	Empty	ABC+*D/

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Priority-Based Stack Operation

- **Left Parenthesis**
 - Behaves as an operator with high **priority** when it is **not** in the stack → **incoming priority (icp) = 0**
 - Behaves as one with **low priority** when it **is in the stack**
→ **in-stack priority (isp) = 8**
 - Only the **matching right parenthesis** can get an in-stack left parenthesis unstacked
- **Summary**
 - Operators are taken out of the stack as long as their **in-stack priority** is **numerically smaller than or equal to the in-coming priority** of the new operator
 - Assuming that the **icp('#') = 8 (lowest)**

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Algorithm of Infix to Postfix

```

void postfix (expression e)
// Output the postfix form of the infix expression e. NextToken
// and stack are as in function eval (Program 3.18). It is assumed that
// the last token in e is '#' Also, '#' is used at the bottom of the stack
{
    Stack<token> stack; // initialize stack
    token y;
    stack.Add('#');
    for (token x = NextToken(e); x != '#'; x = NextToken(e))
    {
        if (x is an operand) cout << x;
        else if (x == ')') // unstack until '('
            for (y = *stack.Delete(y); y != '('; y = *stack.Delete(y)) cout << y;
        else { // x is an operator
            for (y = *stack.Delete(y); isp(y) <= icp(x); y = *stack.Delete(y)) cout << y;
            stack.Add(y); // restack the last y that was unstacked
            stack.Add(x);
        }
    }

    // end of expression; empty stack
    while (!stack.IsEmpty()) cout << *stack.Delete(y);
} // end of postfix

```

higher value
means
lower priority

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Using Stack in STL

C++ reference 網頁: <http://www.cppreference.com/wiki/start>

STL 網頁: <http://www.cppreference.com/wiki/stl/start>

STL Stack 網頁: <http://www.cppreference.com/wiki/stl/stack/start>

Example: the following code uses empty() as the stopping condition on a while loop to clear a stack and display its contents in reverse order:

```

stack<int> s;
for( int i = 0; i < 5; i++ ) { s.push(i); }
while( !s.empty() ) { cout << s.top() << endl; s.pop(); }

```

ch3-54

The End of Chapter 3

**Next Topic:
Linked Lists**